Alternative Approaches to Concurrency

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2

Goals

- Review alternative approaches to concurrency
 - MapReduce
 - Agent Model of Concurrency
 - Examples from Erlang and Scala

Problems with Concurrency

- As mentioned at the beginning of this semester
 - Designing and Implementing Concurrent Systems is hard
 - Data structures shared between threads need to be protected
 - requiring locks (monitors and condition sync) to avoid interference
 - Deadlock and Race conditions are always a concern
 - and sometimes not easy to reproduce
- We've used model-based techniques to attempt to address these concerns; but this simply shifts the problems to the design phase

Alternative Approaches

- As a result of these concerns, computer scientists have searched for other ways to exploit concurrency
 - in particular using techniques from functional programming
- Functional programming is an approach to programming language design in which functions are
 - first class values (with the same status as int or string)
 - you can pass functions as arguments, return them from functions and store them in variables
 - and have no side effects
 - they take input and produce output
 - this typically means that they operate on immutable values

Example (I)

- In python, strings are immutable
 - **a** = "Ken @@@"
 - b = a.replace("@", "!")
 - b
 - \ 'Ken !!!'
 - a
 - 'Ken @@@'
- Replace is a function that takes an immutable value and produces a new immutable value with the desired transformation; it has no side effects

Example (II)

- Functions as values (in python)
 - def Foo(x, y):
 - return x + y
 - add = Foo
 - add(2, 2)
 - **4**
- Here, we defined a function, stored it in a variable, and then used the "call syntax" with that variable to invoke the function that it pointed at

Example (III)

- continuing from previous example
 - def Dolt(fun, x, y): return fun(x,y)
 - Dolt(add, 2, 2)
 - **4**
- Here, we defined a function that accepts three values, some other function and two arguments
 - We then invoked that function by passing our add function along with two arguments;
 - Dolt() is an example of higher-order functions: functions that take functions as parameters
 - ► Higher-order functions is a common idiom in func. prog.

Relationship to Concurrency?

- ► How does this relate to concurrency?
 - It offers a new model for designing concurrent systems
 - Each thread operates on immutable data structures using functions with no side effects
 - A thread's data structures are not shared with other threads
 - Work is performed by passing messages between threads
 - If one thread requires data from another that data is copied and then sent
- Such an approach allows each thread to act like a singlethreaded program; no danger of interference

Map, Filter, Reduce

- Three common higher order functions are map, filter, reduce
- map(fun, list) -> list
 - > Applies fun() to each element of list; returns results in new list
- filter(fun, list) -> list
 - Applies boolean fun() to each element of list; returns new list containing those members of list for which fun() returns True
- reduce(fun, list) -> value
 - Returns a value by applying fun() to successive members of list (total = fun(list[0], list[1]); total = fun(total, list[2]); ...)

Examples

- ightharpoonup list = [10, 20, 30, 40, 50]
- def double(x): return 2 * x
- \triangleright def limit(x): return x > 30
- \triangleright def add(x,y): return x + y
- map(double, list) returns [20, 40, 60, 80, 100]
- filter(limit, list) returns [40, 50]
- reduce(add, list) returns 150

Implications

- map is very powerful
 - especially when you consider that you can pass a list of functions to it and then pass a higher-order function as the function to be applied
 - for example
 - def Dolt(x): return x()
 - map(Dolt, [f(), g(), h(), i(), j(), k()])
- But the real power, with respect to concurrency is that map is simply an abstraction that can, in turn, be implemented in a number of ways

Single Threaded Map

- ► We could for instance implement map() like this:
 - def map(fun, list):
 - results = □
 - for item in list:
 - results.append(fun(item))
- This would implement map in a single threaded fashion

Multi-threaded Map

- ► We could also implement map like this (pseudocode):
 - def Mapper(Thread):
 - def ___init___(... fun, list): ...
 - def run():
 - self.results = map(fun, list)
 - def xmap(fun, list):
 - split list into N parts where N = number of cores
 - create N instances of Mapper(fn, list_i)
 - wait for each thread to end (in order) and grab results
 - append thread results to xmap results
 - return xmap results

Note: threads can complete in any order since each computation is independent

Super Powerful Map

- We could also implement map like this:
 - def supermap(fun, list):
 - divide list into N parts where N equals # of machines
 - send list_i to machine i which then invokes xmap
 - wait for results from each machine
 - combine into single list and return
- Given this implementation, you can apply a very complicated function to a very large list and have (potentially) thousands of machines leap into action to compute the answer

Google

- Indeed, this is what Google does when you submit a search query:
 - \triangleright def aboveThreshold(x): return x > 0.5 <-- just making this up
 - def probabilityDocumentRelatedToSearchTerm(doc): ...
 - searchResults =
 - filter(aboveThreshold,
 - map(probabilityDocumentRelatedToSearchTerm,
 - [<entire contents of the Internet]))</p>

Difference between map and xmap?

- The team behind Erlang published results concerning the difference between map and xmap
 - They make a distinction between
 - CPU-bound computations with little message passing vs.
 - lightweight computations with lots of message passing
- ➤ With the former, xmap provides linear speed-up (10 CPUs provides a 10x speed-up, then declining) over map
 - the latter less so (10 CPUs provided 4x speed-up)
 - Indeed, xmap's performance in the latter case tends to max out at 4x no matter how many CPUs were added

Linear speed-up: Hard to achieve!

- On my machine a program to double each member of a large list actually runs faster in single threaded mode!!
 - When using map, you are building just one results list and do not incur any overhead with respect to threading
 - When using xmap, three lists are being created (one per thread, one to collect the results) and
 - you incur overhead to
 - create each thread
 - wait for each one to start running
 - wait for each one to join the main thread

Demo

Agent Model

- The functional language Erlang is credited with creating an approach to concurrency known as the agent model
 - A concurrent program consists of a set of agents
 - Each agent has its own set of data structures that are not shared with other agents
 - Agents can perform computations and send messages
 - Messages sit in an actor's mailbox until it is ready to process them; they are always processed one at a time
 - An actor does not block when sending a message
 - An actor is not interrupted when a message arrives

Examples

- Examples will be presented in Scala
 - Scala is a language which nicely combines both the imperative and functional programming styles
 - It is implemented on top of Java and thus is cross platform
 - I won't spend much time explaining Scala; I'll just focus on the agent model

Example 1

```
import scala.actors._
object SillyActor extends
  Actor {
  def act() {
    for (i <- 1 to 5) {</pre>
       println("I'm acting!")
       Thread.sleep(1000)
       }
```

```
object SeriousActor extends Actor {
def act() {
for (i <- 1 to 5) {</li>
println("To be or not to be")
Thread.sleep(1000)
}
```

Running Example 1

- SillyActor.start(); SeriousActor.start()
- Demo
 - From this example we can see that Actor is a class that can be sub-classed (just like Thread in Java)
 - You start an actor by calling start()
 - At some point, the scheduler calls the actor's act() method
 - The actor will be active until that method returns
 - This is just like Thread's run() method, only the name has changed

Processing Messages

- To process a message, an actor must call either receive or react
 - react is a special case of receive that we'll discuss below
- You can think of receive as a "switch" statement that specifies the structure of the different type of messages it wants to receive
 - When an actor calls receive, it looks at the mailbox and attempts to find a waiting message that matches one of the branches of the "switch" statement
 - it processes the first match that it finds

Example

receives

This actor loops forever and prints out any message it

Conserve Threads

- When an act() method calls receive(), it tells the scala runtime system that this actor needs its own thread
 - The actor may be spending its time switching between processing messages and performing a long computation
- Since threads in Java are not cheap, scala provides the react keyword to tell the runtime that all this thread does is react to messages
 - This means it spends most of its time blocked
 - Scala uses this information to assign "react actors" to a single thread, thus conserving threads in the overall system

Example

```
object NameResolver extends Actor {
  def act() {
    react {
       case (name: String, actor: Actor) =>
         actor ! getlp(name)
         act()
       case "EXIT" =>
         println("quitting")
```

Note: no explicit loop; that's because react doesn't return (enables sharing of multiple actors on a single thread)

instead, react must call act() if it wants to keep waiting for messages

Results

- To test Scala's claim that react helps conserve threads
 - I wrote a program that can create a specified number of NameResolvers that either
 - use receive or
 - use react
- Results: when creating 100 NameResolvers
 - using receive: 104 threads created
 - using react: 7 threads created (!)

Past Examples

- With the Agent model of concurrency, you can easily avoid interference problems
 - Here's an example of the ornamental garden problem
 - No need for mutual exclusion: create two agents that act as turnstiles and have them send increment messages to a shared counter agent
- However, it can sometimes be tricky to design interactions
 - Here's an example of the museum problem written in this model

Spawning Actors (I)

- The museum example demonstrates a common design idiom in the Agent model of concurrency
 - An agent can only respond to messages when its not doing anything else
 - makes sense: that's just like a single threaded program
 - Think Web browsers and loading images;
 - If they didn't use multiple threads, web pages would load very slowly indeed!
- So, if an agent needs to perform a long computation, it needs to spawn another agent to do that for them

Spawning Agents (II)

```
def reminder() {
  val mainActor = this
  Actor.actor {
    Thread.sleep(1000+generator.nextInt(1000))
    mainActor! "reminder"
  receive {
    case "reminder" =>
      counter! "increment"
      reminder()
```

Final Example

- Message syntax can be as complex as you need it
 - Here's an example of a network node status monitor
 - Taken from this tutorial
 - It queries a domain to see if its "alive"
 - But first
 - Since this example uses Scala "case classes" to create more complex messages with domain-specific syntax
 - Lets do a quick tutorial on case classes

Case Classes

- abstract class Expr
- case class Var(name : String) extends Expr
- case class Number(num: Double) extends Expr
- case class UnOp(operator: String, arg: Expr) extends Expr
- case class BinOp(operator: String, left: Expr, right: Expr) extends Expr
- To create an expression you can now say
- var op = BinOp("+", Number(0), Var(x)) // equals 0 + x

Match Expressions

- Case classes shine when used in match statements
- def simplify(expr : Expr) : Expr = expr match {
 - case UnOp("-", UnOp("-", e)) => e
 - case BinOp("+", Number(0), e) => e
 - case BinOp("*", Number(1), e) => e
 - case _ => expr
- simplify(BinOp("+", Number(0), Var(x))) returns Var(x)
 - ightharpoonup because "0 + x" == "x"

Case classes in Example

- case class NodeStatusRequest(address: InetAddress, a: Actor)
- sealed abstract class NodeStatus
- case class Available(address: InetAddress) extends NodeStatus
- case class Unresponsive(
 - address: InetAddress,
 - reason: Option[String]) extends NodeStatus
- Option[String] is a special type that can either have the value Some(String) or None

Demo

- This demo sets up an actor to check the availability of various domains
- It then passes a few messages to this actor and then waits for the actor to respond
 - It can also handle the case when it gets an unexpected message

Wrapping Up

- We have looked at a few alternative models to the "locks and shared data" model of concurrency that
 - draw on functional programming techniques
 - do not allow threads to share data
 - allow threads to communicate via asynchronous messages
- Deadlock and Race conditions are still possible in this model but harder to achieve
 - However, interference is simply not possible in this model
- Functional techniques seem like a promising method for tackling concurrency on multi-core hardware

Coming Up

- Lecture 27: Dealing with Bugs
 - Chapter 11 of Head First Software Development
- Lecture 28: Software Abstractions
 - Overview of Software Abstractions Optional Textbook